



# Introduction to Cryptography

## Lecture 6

Benny Pinkas





# Public Key Encryption



# Classical symmetric ciphers

- Alice and Bob share a private key  $k$ .
- System is secure as long as  $k$  is secret.
- Major problem: generating and distributing  $k$ .





## Diffie and Hellman: “New Directions in Cryptography”, 1976.

- “We stand today on the brink of a revolution in cryptography. The development of cheap digital hardware has freed it from the design limitations of mechanical computing...  
...such applications create a need for new types of cryptographic systems which minimize the necessity of secure key distribution...  
...theoretical developments in information theory and computer science show promise of providing provably secure cryptosystems, changing this ancient art into a science.”



# Diffie-Hellman

- Came up with the idea of public key cryptography



Everyone can learn Bob's public key and encrypt messages to Bob.  
Only Bob knows the decryption key and can decrypt.

Key distribution is greatly simplified.



# Plan

- Basic number theory
  - Divisors, modular arithmetic
  - The GCD algorithm
  - Groups
- References:
  - Many books on number theory
  - Almost all books on cryptography
  - Cormen, Leiserson, Rivest, (Stein), “Introduction to Algorithms”, chapter on Number-Theoretic Algorithms.



# Divisors, prime numbers

- We work over the integers
- A non-zero integer  $b$  divides an integer  $a$  if there exists an integer  $c$  s.t.  $a=c\cdot b$ .
  - Denoted as  $b|a$
  - I.e.  $b$  divides  $a$  with no remainder
- Examples
  - Trivial divisors:  $1|a$ ,  $a|a$
  - Each of  $\{1,2,3,4,6,8,12,24\}$  divides 24
  - 5 does not divide 24
- Prime numbers
  - An integer  $a$  is prime if it is only divisible by 1 and by itself.
  - 23 is prime, 24 is not.



# Modular Arithmetic

- Modular operator:
  - $a \bmod b$ , (or  $a \% b$ ) is the remainder of  $a$  when divided by  $b$
  - I.e., the smallest  $r \geq 0$  s.t.  $\exists$  integer  $q$  for which  $a = qb + r$ .
  - (Thm: there is a single choice for such  $q, r$ )
- Examples
  - $12 \bmod 5 = 2$
  - $10 \bmod 5 = 0$
  - $-5 \bmod 5 = 0$
  - $-1 \bmod 5 = 4$



# Modular congruency

- $a$  is congruent to  $b$  modulo  $n$  ( $a \equiv b \pmod{n}$ ) if
  - $(a-b) = 0 \pmod{n}$
  - Namely,  $n$  divides  $a-b$
  - In other words,  $(a \pmod{n}) = (b \pmod{n})$
- E.g.,
  - $23 \equiv 12 \pmod{11}$
  - $4 \equiv -1 \pmod{5}$
- There are  $n$  equivalence classes modulo  $n$ 
  - $[3]_7 = \{\dots, -11, -4, 3, 10, 17, \dots\}$



# Greatest Common Divisor (GCD)

- $d$  is a common divisor of  $a$  and  $b$ , if  $d|a$  and  $d|b$ .
- $\gcd(a,b)$  (Greatest Common Divisor), is the largest integer that divides both  $a$  and  $b$ . ( $a,b \geq 0$ )
  - $\gcd(a,b) = \max k$  s.t.  $k|a$  and  $k|b$ .
- Examples:
  - $\gcd(30,24) = 6$
  - $\gcd(30,23) = 1$
- If  $\gcd(a,b)=1$  then  $a$  and  $b$  are denoted relatively prime.



# Facts about the GCD

- $\gcd(a,b) = \gcd(b, a \bmod b)$  (interesting when  $a > b$ )
- Since (e.g.,  $a=33, b=15$ )
  - If  $c|a$  and  $c|b$  then  $c|(a \bmod b)$
  - If  $c|b$  and  $c|(a \bmod b)$  then  $c|a$
- If  $a \bmod b = 0$ , then  $\gcd(a,b)=b$ .
- Therefore,

$$\begin{aligned}\gcd(19,8) &= \\ \gcd(8, 3) &= \\ \gcd(3,2) &= \\ \gcd(2,1) &= 1\end{aligned}$$

$$\begin{aligned}\gcd(20,8) &= \\ \gcd(8, 4) &= 4\end{aligned}$$



# Euclid's algorithm

Input:  $a > b > 0$

Output:  $\gcd(a, b)$

Algorithm:

1. if  $(a \bmod b) = 0$  return  $(b)$
2. else return  $(\gcd(b, a \bmod b))$

Complexity:

- $O(\log a)$  rounds
- Each round requires  $O(\log^2 a)$  bit operations
- Actually, the total overhead can be shown to be  $O(\log^2 a)$



# The extended gcd algorithm

Finding  $s, t$  such that  $\gcd(a,b) = a \cdot s + b \cdot t$

Extended-gcd(a,b) /\* output is (gcd(a,b), s, t)

1. If  $(a \bmod b = 0)$  then return(b,0,1)
2.  $(d',s',t') = \text{Extended-gcd}(b, a \bmod b)$
3.  $(d,s,t) = (d', t', s' - \lfloor a/b \rfloor \cdot t')$
4. return(d,s,t)

Note that the overhead is as in the basic GCD algorithm



- Extended gcd algorithm
  - Given  $a, b$  finds  $s, t$  such that  $\gcd(a, b) = a \cdot s + b \cdot t$
  - In particular, if  $p$  is prime and  $a < p$  then  $\gcd(a, p) = 1$ , and therefore  $a \cdot s + p \cdot t = 1$ .



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- THM: There is no positive integer smaller than  $\gcd(a, b)$  which can be represented as a linear combination of  $a, b$ .
  - For example,  $a=12, b=8$ .
  - $4 = 1 \cdot 12 - 1 \cdot 8$
  - There are no  $s, t$  for which  $2 = s \cdot 12 + t \cdot 8$
- Therefore if we find  $s, t$  such that  $as + tb = 1$ , then we know that  $\gcd(a, b) = 1$



# Groups

- Definition: a set  $G$  with a binary operation  $\circ: G \times G \rightarrow G$  is called a group if:
  - (closure)  $\forall a, b \in G$ , it holds that  $a \circ b \in G$ .
  - (associativity)  $\forall a, b, c \in G$ ,  $(a \circ b) \circ c = a \circ (b \circ c)$ .
  - (identity element)  $\exists e \in G$ , s.t.  $\forall a \in G$  it holds that  $a \circ e = a$ .
  - (inverse element)  $\forall a \in G \exists a^{-1} \in G$ , s.t.  $a \circ a^{-1} = e$ .
- A group is Abelian (commutative) if  $\forall a, b \in G$ , it holds that  $a \circ b = b \circ a$ .
- Examples:
  - Integers under addition
    - $(\mathbb{Z}, +) = \{\dots, -3, -2, -1, 0, 1, 2, 3, \dots\}$



# More examples of groups

- Addition modulo N
  - $(G, \circ) = (\{0, 1, 2, \dots, N-1\}, +)$



# More examples of groups

- $Z_p^*$  Multiplication modulo a prime number  $p$ 
  - $(G, \circ) = (\{1, 2, \dots, p-1\}, \times)$
  - E.g.,  $Z_7^* = (\{1, 2, 3, 4, 5, 6\}, \times)$
- Closure (the result of the multiplication is never divisible by  $p$ ):
  - $s \cdot a + t \cdot p = 1$
  - $s' \cdot b + t' \cdot p = 1$
  - $ss' \cdot (ab) + (sat' + s'bt + tt'p) \cdot p = 1$
  - Therefore  $1 = \gcd(ab, p)$ .
- Trivial: associativity, existence of identity element.
- The extended GCD algorithm shows that an inverse always exists:  $s \cdot a + t \cdot p = 1 \Rightarrow s \cdot a = 1 - t \cdot p \Rightarrow s \cdot a \equiv 1 \pmod{p}$



# More examples of groups

- $Z_N^*$  Multiplication modulo a composite number  $N$ 
  - $(G, \circ) = (\{a \text{ s.t. } 1 \leq a \leq N-1 \text{ and } \gcd(a, N)=1\}, \times)$
  - E.g.,  $Z_{10}^* = (\{1, 3, 7, 9\}, \times)$
  - Closure:
    - $s \cdot a + t \cdot N = 1$
    - $s' \cdot b + t' \cdot N = 1$
    - $ss' \cdot (ab) + (sat' + s'bt + tt'N) \cdot N = 1$
    - Therefore  $1 = \gcd(ab, N)$ .
  - Associativity: trivial
  - Existence of identity element: 1.
  - Inverse element: as in  $Z_p^*$



# Subgroups

- Let  $(G, \circ)$  be a group.
  - $(H, \circ)$  is a subgroup of  $G$  if
    - $(H, \circ)$  is a group
    - $H \subseteq G$
  - For example,  $H = (\{1, 2, 4\}, \times)$  is a subgroup of  $Z_7^*$ .
- *Lagrange's theorem:*  
If  $(G, \circ)$  is finite and  $(H, \circ)$  is a subgroup of  $(G, \circ)$ , then  $|H|$  divides  $|G|$

In our example:  $3|6$ .



# Cyclic Groups

- Exponentiation is repeated application of  $\circ$ 
  - $a^3 = a \circ a \circ a$ .
  - $a^0 = 1$ .
  - $a^{-x} = (a^{-1})^x$
- A group  $G$  is cyclic if there exists a generator  $g$ , s.t.  
 $\forall a \in G, \exists i$  s.t.  $g^i = a$ .
  - I.e.,  $G = \langle g \rangle = \{1, g, g^2, g^3, \dots\}$
  - For example  $Z_7^* = \langle 3 \rangle = \{1, 3, 2, 6, 4, 5\}$
- Not all  $a \in G$  are generators of  $G$ , but they all generate a subgroup of  $G$ .
  - E.g. 2 is not a generator of  $Z_7^*$
- The order of a group element  $a$  is the smallest  $j > 0$  s.t.  
 $a^j = 1$
- *Lagrange's theorem*  $\Rightarrow$  for  $x \in Z_p^*$ ,  $\text{ord}(x) \mid p-1$ .



# Fermat's theorem

- Corollary of Lagrange's theorem: if  $(G, \circ)$  is a finite group, then  $\forall a \in G, a^{|G|} = 1$ .
- Proof:
  - $a \in G$  generates a subgroup  $H$  of  $G$ .
  - $A^{|H|} = 1$
  - Lagrange theorem:  $|H| \mid |G|$
  - $A^{|G|} = (A^{|H|})^{|G|/|H|} = 1$



# Fermat's theorem

- Corollary of Lagrange's theorem: if  $(G, \circ)$  is a finite group, then  $\forall a \in G, a^{|G|} = 1$ .
- Corollary (Fermat's theorem):  $\forall a \in \mathbb{Z}_p^*, a^{p-1} = 1 \pmod p$ .  
E.g., for all  $\forall a \in \mathbb{Z}_7^*, a^6 = 1, a^7 = a$ .
- Computing inverses:
- Given  $a \in G$ , how to compute  $a^{-1}$ ?
  - Fermat's theorem:  $a^{-1} = a^{|G|-1}$  ( $= a^{p-2}$  in  $\mathbb{Z}_p^*$ )
  - Or, using the extended gcd algorithm (for  $\mathbb{Z}_p^*$  or  $\mathbb{Z}_N^*$ ):
    - $\gcd(a, p) = 1$
    - $s \cdot a + t \cdot p = 1 \Rightarrow s \cdot a = -t \cdot p + 1 \Rightarrow s$  is  $a^{-1}$  !!
  - Which is more efficient?



# Computing in $Z_p^*$

- P is a huge prime (1024 bits)
- Easy tasks (measured in bit operations):
  - Adding in  $O(\log p)$  (namely, linear in the length of p)
  - Multiplying in  $O(\log^2 p)$  (and even in  $O(\log^{1.7} p)$ )
  - Inverting ( $a$  to  $a^{-1}$ ) using extended gcd in  $O(\log^2 p)$
  - Exponentiations  $x^r \bmod p$  in  $O(\log r \cdot \log^2 p)$ , using repeated squaring:
    - $r = r_{|r|-1} \dots r_1 x_0$ . Therefore  $x^r = x^{\sum_{i=0}^{|r|-1} 2^i r_i} = \prod_{i=0}^{|r|-1} x^{2^i r_i} = \prod_{r_i=1} x^{2^i}$
    - Algorithm:
      - out = 1; b=x
      - For i = 0 to |r|-1
        - If  $r_i = 1$  then out = out \* b
        - b = b \* b



# Groups we will use

- $Z_p^*$  Multiplication modulo a prime number  $p$ 
  - $(G, \circ) = (\{1, 2, \dots, p-1\}, \times)$
  - E.g.,  $Z_7^* = (\{1, 2, 3, 4, 5, 6\}, \times)$
- $Z_N^*$  Multiplication modulo a composite number  $N$ 
  - $(G, \circ) = (\{a \text{ s.t. } 1 \leq a \leq N-1 \text{ and } \gcd(a, N)=1\}, \times)$
  - E.g.,  $Z_{10}^* = (\{1, 3, 7, 9\}, \times)$



# Euler's phi function

- Lagrange's Theorem:  $\forall a$  in a finite group  $G$ ,  $a^{|G|}=1$ .
- Euler's phi function (aka, Euler's totient function),
  - $\phi(n)$  = number of elements in  $Z_n^*$  (i.e.  $|\{x \mid \gcd(x,n)=1, 1 \leq x \leq n\}|$ )
  - $\phi(p) = p-1$  for a prime  $p$ .
  - $n = \prod_{i=1..k} p_i^{e(i)} \Rightarrow \phi(n) = n \cdot \prod_{i=1..k} (1 - 1/p_i)$
  - $\phi(p^2) = p(p-1)$  for a prime  $p$ .
  - $n = p \cdot q \Rightarrow \phi(n) = (p-1)(q-1)$
- Corollary: For  $Z_n^*$  ( $n = p \cdot q$ ),  $|Z_n^*| = \phi(n) = (p-1)(q-1)$ .
- $\forall a \in Z_n^*$  it holds that  $a^{\phi(n)} = 1 \pmod n$ 
  - For  $Z_p^*$  (prime  $p$ ),  $a^{p-1} = 1 \pmod p$  (Fermat's theorem).
  - For  $Z_n^*$  ( $n = p \cdot q$ ),  $a^{(p-1)(q-1)} = 1 \pmod n$





# Finding prime numbers



# Finding prime numbers

- Prime number theorem:  $\#\{\text{primes} \leq x\} \approx x / \ln x$  as  $x \rightarrow \infty$
- How can we find a random  $k$ -bit prime?
  - Choose  $x$  at random in  $\{2^k, \dots, 2^{k+1}-1\}$ 
    - (How many numbers in that range are prime?  
About  $2^{k+1} / \ln 2^{k+1} - 2^k / \ln 2^k$  numbers, i.e.  $\approx$  a  $1 / \ln(2^k)$  fraction.)
  - Test if  $x$  is prime
    - (more on this later in the course)
- The probability of success is  $\approx 1 / \ln(2^k) = O(1/k)$ .
- The expected number of trials is  $O(k)$ .



# Finding generators

- How can we find a generator of  $Z_p^*$ ?
- Pick a random number  $a \in [1, p-1]$ , check if is a generator
  - Naively, check whether  $\forall 1 \leq i \leq p-2 \quad a^i \neq 1$  ☹
  - But we know that if  $a^i = 1 \pmod p$  then  $i \mid p-1$ .
  - Therefore need to only check  $i$  for which  $i \mid p-1$ .
- Easy if we know the factorization of  $(p-1)$ . In that case
  - For all  $a \in Z_p^*$ , the order of  $a$  divides  $(p-1)$
  - For every integer divisor  $b$  of  $(p-1)$ , check if  $a^b = 1 \pmod p$ .
  - If none of these checks succeeds, then  $a$  is a generator, since its order must be  $p-1$ .



# Finding prime numbers of the right form

- How can we know the factorization of  $p-1$ ?
  - Easy, for example, if  $p=2q+1$ , and  $q$  is prime.
  - How can we find a  $k$ -bit prime of this form?
1. Search for a prime number  $q$  of length  $k-1$  bits. (Will be successful after about  $O(k)$  attempts.)
  2. Check if  $2q+1$  is prime (we will see how to do this later in the course).
  3. If not, go to step 1.