Introduction to Cryptography: *Homework 2*

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1. Let p be a prime number such that $p\text{-}1=p_1^{e1}p_2^{e2}\ldots p_m^{em}$ ($\forall i$, p_i is prime and $e_i \geq 1$). Prove that $g \in Z_p^*$ is a generator if and only if for all $1 \leq i \leq m$ it holds that $g^{(p-1)/pi} \neq 1 \mod p$.

2. The purpose of this exercise is to find an efficient algorithm for computing discrete logarithms in Z_p^*, where p is prime and $p=2^n+1$.
   The discrete logarithm problem is the following:
   Input: a prime p, a generator g of Z_p^*, and a value y in Z_p^*.
   Output: x s.t. $g^x = y \mod p$.

   Let $x=b_{n-1}2^{n-1}+ b_{n-2}2^{n-2}+\ldots+b_1 2^1+b_0$ be the binary representation of x.

   a. Show how to find the least significant bit (b_0) of x (given g,y). (7 points)
   b. Set $z=y \cdot g^{-b0}$, and show how to use it to find the bit b_1. (10 points)
      Hint: there is an integer i such that $z=g^{4i+2\cdot b1}$. Recall also that $e=p\text{-}1=2^n$ is the smallest exponent s.t. $g^e=1 \mod p$. Use these facts to find b_1.
   c. Show how to find the complete binary representation of x. (10 points)
   d. Explain why this method is only good for a prime modulo p that satisfies $p=2^n+1$. (6 points)

   Note: this algorithm can be generalized for any Z_p^* for which $p\text{-}1=p_1^{e1}p_2^{e2}\ldots p_m^{em}$, all p_i are small primes, and the factorization of p-1 is known. (There is not need to prove this fact.)

3. Let p be a prime number. Suppose that g is a generator of Z_p^* and let $b=g^i$ for an exponent $0 \leq i \leq p\text{-}2$.
   a. Show that the order of b is $(p\text{-}1)/gcd(p\text{-}1,i)$. (17 points)
   b. Show that the number of generators in Z_p^* is ϕ(p-1). (16 points)